

Boolean Algebra Logic Gates

Introduction

- Hardware consists of a few simple building blocks
 - These are called *logic gates*
 - AND, OR, NOT, ...
 - NAND, NOR, XOR, ...
- Logic gates are built using transistors
 - NOT gate can be implemented by a single transistor
 - AND gate requires 3 transistors
- Transistors are the fundamental devices
 - Pentium processor consists of 3 million transistors
 - Now we can build chips with more than 100 million transistors

Basic Concepts

- Simple gates
 - AND
 - OR
 - NOT
- Functionality can be expressed by a truth table
 - A truth table lists output for each possible input combination
- Precedence
 - NOT > AND > OR
 - $F = A \overline{B} + \overline{A} B$ $= (A (\overline{B})) + ((\overline{A}) B)$



- Additional useful gates
 - NAND
 - NOR
 - XOR
- NAND = AND + NOT
- NOR = OR + NOT
- XOR implements exclusive-OR function
- NAND and NOR gates require only 2 transistors
 - AND and OR need 3 transistors!



- Number of functions
 - With N logical variables, we can define 2^{2^N} functions
 - Some of them are useful
 - AND, NAND, NOR, XOR, ...
 - Some are not useful:
 - Output is always 1
 - Output is always 0
 - "Number of functions" definition is useful in proving completeness property

- Complete sets
 - A set of gates is complete
 - If we can implement any logical function using only the type of gates in the set
 - You can uses as many gates as you want
 - Some example complete sets
 - {AND, OR, NOT} Not a minimal complete set
 - {AND, NOT}
 - {OR, NOT}
 - {NAND}
 - {NOR}

- Minimal complete set

A complete set with no redundant elements.

• Proving NAND gate is universal



• Proving NOR gate is universal



Logic Chips







Logic Chips (cont.)

- Integration levels
 - SSI (small scale integration)
 - Introduced in late 1960s
 - 1-10 gates (previous examples)
 - MSI (medium scale integration)
 - Introduced in late 1960s
 - 10-100 gates

- LSI (large scale integration)

- Introduced in early 1970s
- 100-10,000 gates

- VLSI (very large scale integration)

- Introduced in late 1970s
- More than 10,000 gates

Logic Functions

- Logical functions can be expressed in several ways:
 - Truth table
 - Logical expressions
 - Graphical form

• Example:

- Majority function
 - Output is one whenever majority of inputs is 1
 - We use 3-input majority function

Logic Functions (cont.)

3-input majority function

Α	В	С	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

Logical expression form
F = A B + B C + A C



Logical Equivalence

• All three circuits implement F = A B function



Logical Equivalence (cont.)

- Proving logical equivalence of two circuits
 - Derive the logical expression for the output of each circuit
 - Show that these two expressions are equivalent
 - Two ways:
 - You can use the truth table method
 - » For every combination of inputs, if both expressions yield the same output, they are equivalent
 - » Good for logical expressions with small number of variables
 - You can also use algebraic manipulation
 - » Need Boolean identities

Logical Equivalence (cont.)

- Derivation of logical expression from a circuit
 - Trace from the input to output
 - Write down intermediate logical expressions along the path



Logical Equivalence (cont.)

• Proving logical equivalence: Truth table method

Α	В	F1 = A B	$F3 = (A + B) (A + \overline{B}) (\overline{A} + B)$
0	0	0	0
0	1	0	0
1	0	0	0
1	1	1	1

Boolean Algebra

Boolean identities

Name	AND version	OR version
Identity	$x \cdot 1 = x$	$\mathbf{x} + 0 = \mathbf{x}$
Complement	$\mathbf{x} \cdot \overline{\mathbf{x}} = 0$	$x + \overline{x} = 1$
Commutative	$\mathbf{x} \cdot \mathbf{y} = \mathbf{y} \cdot \mathbf{x}$	$\mathbf{x} + \mathbf{y} = \mathbf{y} + \mathbf{x}$
Distribution	x(y+z) = xy+xz	$\mathbf{x} + (\mathbf{y} \cdot \mathbf{z}) =$
		(x+y) (x+z)
Idempotent	$\mathbf{X} \cdot \mathbf{X} = \mathbf{X}$	$\mathbf{x} + \mathbf{x} = \mathbf{x}$
Null	$\mathbf{X} 0 = 0$	x + 1 = 1

Boolean Algebra (cont.)

Name	AND version	OR version
Involution	$\overline{\overline{\mathbf{x}}} = \mathbf{x}$	
Absorption	x (x+y) = x	$\mathbf{x} + (\mathbf{x} \cdot \mathbf{y}) = \mathbf{x}$
Associative	$\mathbf{x} \cdot (\mathbf{y} \cdot \mathbf{z}) = (\mathbf{x} \cdot \mathbf{y}) \cdot \mathbf{z}$	$\mathbf{x} + (\mathbf{y} + \mathbf{z}) =$
		(x+y)+z
de Morgan	$\overline{\mathbf{x}} \cdot \mathbf{y} = \overline{\mathbf{x}} + \overline{\mathbf{y}}$	$\overline{\mathbf{x} + \mathbf{y}} = \overline{\mathbf{x}} \cdot \overline{\mathbf{y}}$

Boolean Algebra (cont.)

- Proving logical equivalence: Boolean algebra method
 - To prove that two logical functions F1 and F2 are equivalent
 - Start with one function and apply Boolean laws to derive the other function
 - Needs intuition as to which laws should be applied and when
 - Practice helps!
 - Sometimes it may be convenient to reduce both functions to the same expression
 - Example: F1= A B and F3 are equivalent

Logic Circuit Design Process

- A simple logic design process involves
 - Problem specification
 - Truth table derivation
 - Derivation of logical expression
 - Simplification of logical expression
 - Implementation



Deriving Logical Expressions

- Derivation of logical expressions from truth tables
 - sum-of-products (SOP) form
 - product-of-sums (POS) form
- SOP form
 - Write an AND term for each input combination that produces a 1 output
 - Write the variable if its value is 1; complement otherwise
 - OR the AND terms to get the final expression
- POS form
 - Dual of the SOP form

Deriving Logical Expressions

• 3-input majority function

Α	В	С	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

- SOP logical expression
- Four product terms
 - Because there are 4 rows with a 1 output

$$F = \overline{A} B C + A \overline{B} C + A \overline{B} C + A B \overline{C} + A B C$$

Deriving Logical Expressions

• 3-input majority function

Α	В	С	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

- POS logical expression
- Four sum terms
 - Because there are 4 rows with a 0 output

$$F = (A + B + C) (A + B + C)$$

(A + B + C) (A + B + C)
(A + B + C) (A + B + C)

Logical Expression Simplification

- -Algebraic manipulation
 - Use Boolean laws to simplify the expression
 - Difficult to use
 - Don't know if you have the simplified form

Algebraic Manipulation

- Majority function example Added extra $\overline{A}BC + A\overline{B}C + AB\overline{C} + ABC =$ $\overline{A}BC + A\overline{B}C + AB\overline{C} + ABC + ABC + ABC$ We can now simplify this expression as BC + AC + AB
- A difficult method to use for complex expressions

Implementation Using NAND Gates

- Using NAND gates
 - Get an equivalent expression

A B + C D = A B + C D

- Using de Morgan's law

 $A B + C D = A B \cdot C D$

- Can be generalized
 - Majority function

$$A B + B C + AC = \overline{A B} \cdot \overline{BC} \cdot \overline{AC}$$

Idea: NAND Gates: Sum-of-Products, NOR Gates: Product-of-Sums

Implementation Using NAND Gates (cont.)

Majority function



Introduction to Combinational Circuits

- Combinational circuits
 - Output depends only on the current inputs
- Combinational circuits provide a higher level of abstraction
 - Help in reducing design complexity
 - Reduce chip count
- We look at some useful combinational circuits