ACTIVE REDUNDANCY

When the reliability of a series system does not reach the design goal, it becomes necessary to act at the structure level and to resort to <u>redundant</u> configurations.

A system configuration is said to be <u>redundant</u>, when the occurrence of <u>a component failure does not necessarily causes a</u> <u>system failure</u>.

Various redundant architectures have been studied and applied in practice, and they will be illustrated in the following sections.

From the operating point of view, we can distinguish between:

- Active Redundancy (parallel, hot): Redundant elements are subjected from the beginning to <u>the same load as the</u> <u>operating elements</u>;
- Warm Redundancy (lightly loaded): Redundant elements are subjected to <u>a lower load</u> until they become operating;
- *Standby Redundancy (cold, unloaded)*: Redundant elements are subjected to <u>no load</u> until they become operating, and the failure rate in reserve (standby) state is <u>assumed to be zero</u>.

By default, we assume the <u>components</u> of the system with redundancy <u>are statistically independent</u>, i.e. the failure of one of them doesn't affect the reliability of the rest.

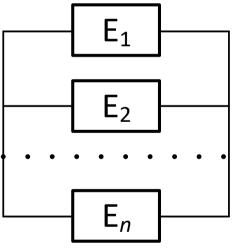
However, if this is not the case, i.e. the component(s) that are still operating assume the failed unit's portion of the load, such type of redundancy is called *load sharing*.

The reliability of load sharing configurations is much harder to compute.

A <u>parallel model</u> consists of n (<u>often statistically identical</u>) elements in active redundancy, of which k ($1 \le k < n$) are necessary to perform the required function and the remaining (n - k) are in reserve.

Such a structure is designated as a *k-out-of-n* (or *k-out-of-n:G*) redundancy.

1-out-of-n redundancy is also called <u>hot</u> <u>redundancy</u>.

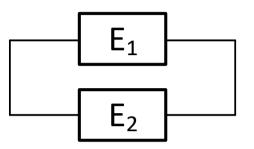


Let's consider at first the case of an active (hot) *1-out-of-2* redundancy.

The required function is fulfilled if at least one of the elements E_1 or E_2 works without failure in the interval (0, t]. In other words, the system fails if both elements failed in the interval (0, t].

Let $\{\overline{e_i}\}, i = 1, 2$ be the event of *i*th element failing in (0, t], then $F_i(t) = Pr\{\overline{e_i}\}$ is the failure probability of E_i .

Then, if $F_s(t)$ denotes the failure probability of the entire system, it follows:



$$F_{S}(t) = Pr\{\overline{e_{1}} \cap \overline{e_{2}}\} = Pr\{\overline{e_{1}}\} \cdot Pr\{\overline{e_{2}}\} = F_{1}(t) \cdot F_{2}(t)$$

Generalizing for 1-out-of-*n* system, we obtain

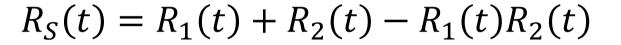
$$F_S(t) = \prod_{i=1}^n F_i(t)$$

Product Law of Unreliabilities

 E_1

 E_2

Substituting $F_*(t)$ with $1 - R_*(t)$, we obtain



For *1-out-of-n* system, we get

$$R_{S}(t) = 1 - \prod_{i=1}^{n} (1 - R_{i}(t))$$

As it is often that all elements in parallel system are statistically identical, i.e. $R_1(t) = R_2(t) = \cdots = R_n(t) = R(t)$, we can obtain for this particular case

$$R_S(t) = 2R(t) - R(t)^2$$
 for 1-out-of-2 system

and

$$R_S(t) = 1 - (1 - R(t))^n$$
 for 1-out-of-n system.

Thus, for parallel systems of independent components, we have a <u>product law of "unreliabilities"</u> analogous to the product law of reliabilities for series systems.

It follows that the reliability of a parallel system is greater than the reliability of each constituent component ...

... and, hence, a parallel system is more reliable than the most reliable of its components.

Let's assume the parallel system consists of 2 identical components, and

$$R_1(t) = R_2(t) = R(t) = e^{-\lambda t},$$

i.e. the failure time of each component is an exponentially distributed r.v.

Moreover, the failure rate of each component is constant.

The reliability of the system then is

$$R_S(t) = 2R(t) - R(t)^2 = 2e^{-\lambda t} - e^{-2\lambda t}$$

Obviously, we can't find such $\lambda_s > 0$ so that

$$e^{-\lambda_S t} = 2e^{-\lambda t} - e^{-2\lambda t}$$

This means that <u>the failure rate of the entire system is not</u> <u>constant</u>, and its failure time is <u>not an exponentially distributed</u> <u>*r.v.*</u>

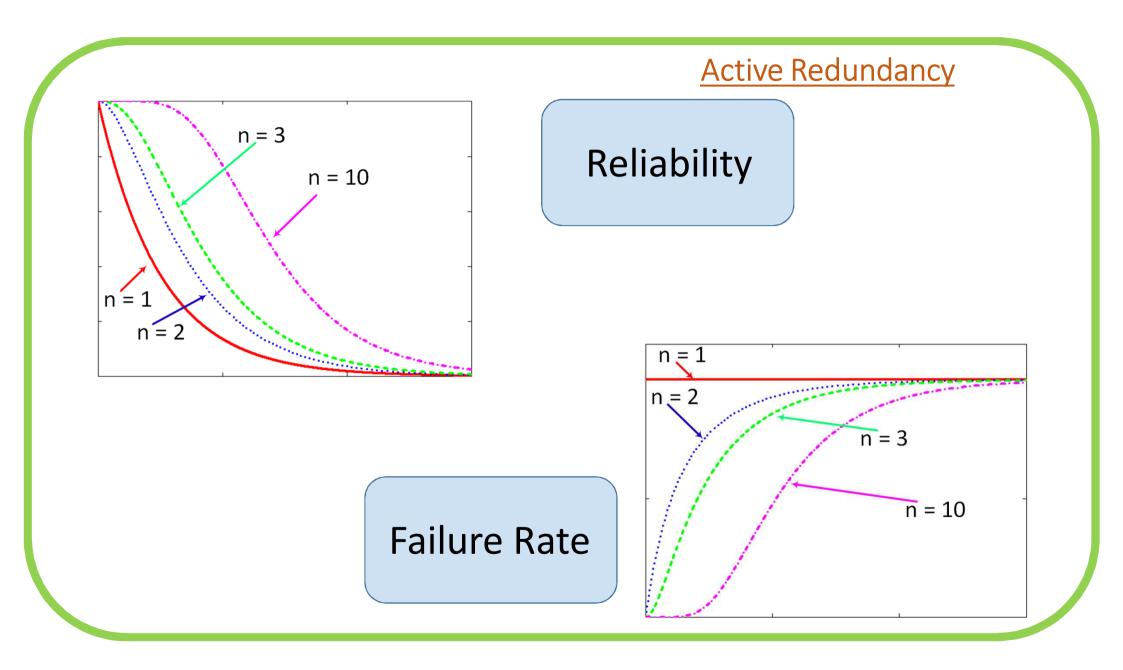
To demonstrate this, let's find the expression for the system's failure rate.

We know that
$$h(t) = -\frac{R'(t)}{R(t)}$$

$$R_S(t) = 2e^{-\lambda t} - e^{-2\lambda t}$$

we obtain

$$h_{S}(t) = \frac{2\lambda e^{-\lambda t} - 2\lambda e^{-2\lambda t}}{2e^{-\lambda t} - e^{-2\lambda t}} = \frac{2\lambda e^{-\lambda t} \left(1 - e^{-\lambda t}\right)}{e^{-\lambda t} \left(2 - e^{-\lambda t}\right)} = \frac{2\lambda \left(1 - e^{-\lambda t}\right)}{2 - e^{-\lambda t}}.$$



As previously mentioned, a hot redundancy system is the special case of *k-out-of-n* redundancy systems (with *k=1*).

Mind that for *k-out-of-n* systems correct operation of *k* components is sufficient for the system to be operable, hence the system can tolerate (*n-k*) failures of its components.

By default we assume that all items in *k-out-of-n* systems are statistically identical.

Consider a *k*-out-of-n system $(2 \le k < n)$ with reliability of each item equals *R*.

The system is operable if it has *n*, *n*-1, *n*-2, ..., *k*+1, *k* operable components.

The probability that exactly n-1 components are operable is given by $n \cdot R^{n-1} \cdot (1-R)^1$. Here the multiplier n is stands for the number of all possible combinations of operating items, i.e.

$$\binom{n}{n-1} = \frac{n!}{(n-1)!(n-n+1)!} = n.$$

The probability that exactly *n*-2 components are operable is given by $\binom{n}{n-2} \cdot R^{n-2} \cdot (1-R)^2$.

Finally, the probability that exactly k components are operable is given by $\binom{n}{k} \cdot R^k \cdot (1-R)^{n-k}$.

. . .

Since we have listed all possible conditions for the system to be operable, the reliability of the system is given by

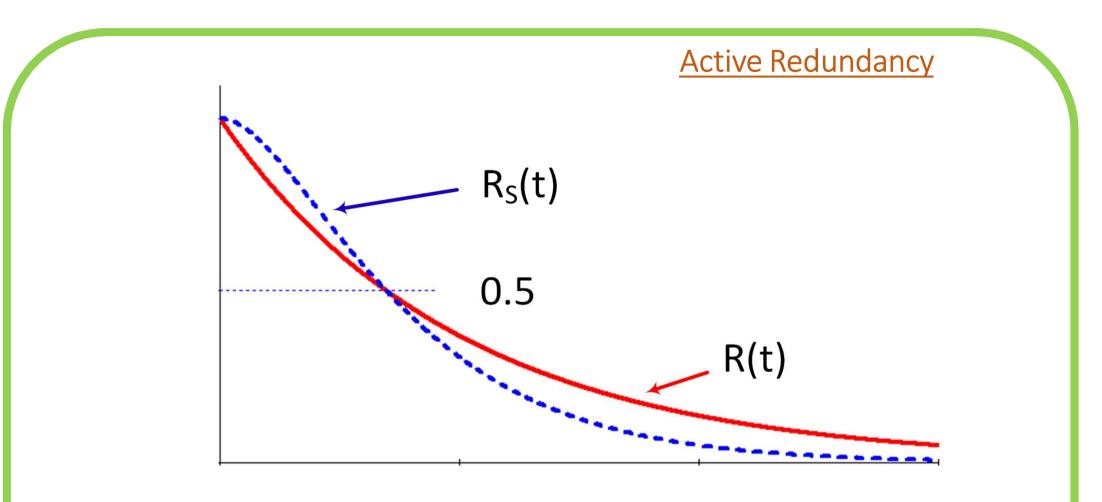
$$R_S = \sum_{i=k}^n \binom{n}{i} \cdot R^i \cdot (1-R)^{n-i}$$

Ex.: Consider a 2-out-of-3 system with reliability of each component given as $R(t) = e^{-\lambda t}$. Find the reliability function, MTTF and failure rate function for the entire system.

First, the reliability of the system is obtained as

$$R_{S}(t) = \sum_{i=2}^{3} {\binom{3}{i}} \cdot R(t)^{i} \cdot (1 - R(t))^{3-i} =$$

= $3R(t)^{2}(1 - R(t)) + R(t)^{3} =$
= $3R(t)^{2} - 2R(t)^{3} =$
= $3e^{-2\lambda t} - 2e^{-3\lambda t}$



The system is more reliable if the reliability of each item is greater than 0,5.

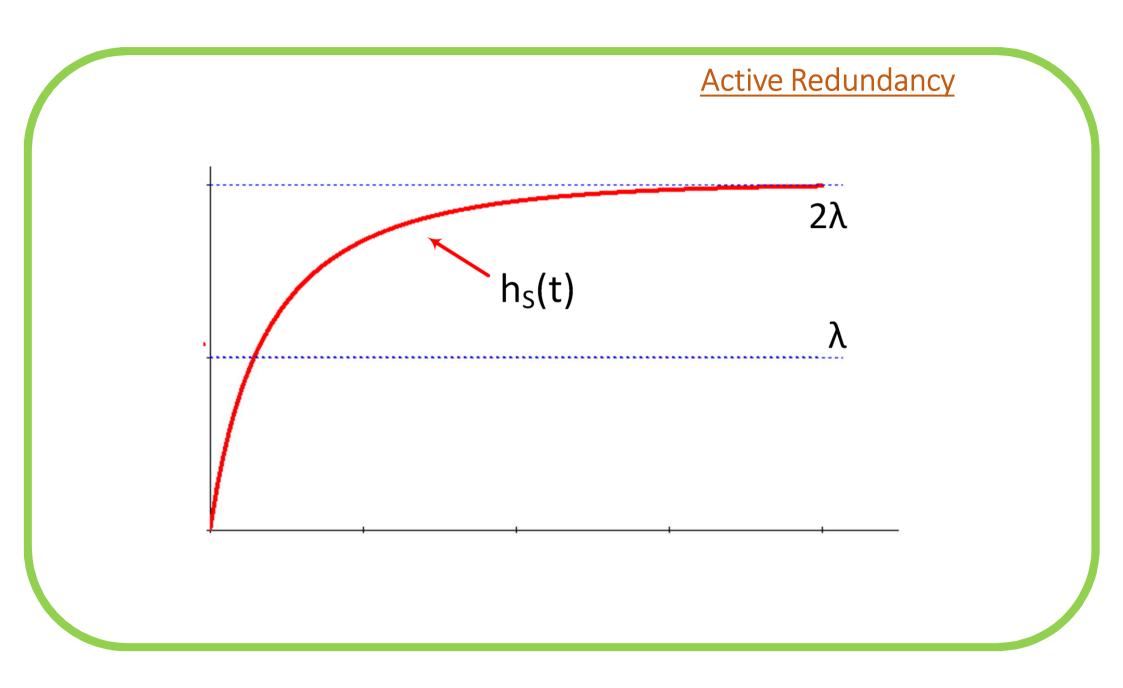
Second, the MTTF is given by

$$MTTF = \int_{0}^{\infty} R_{S}(t)dt = \int_{0}^{\infty} \left(3e^{-2\lambda t} - 2e^{-3\lambda t}\right)dt =$$
$$= \frac{3}{2\lambda} - \frac{2}{3\lambda} = \frac{5}{6\lambda}.$$

Note that MTTF of the *2-out-of-3* system is less than the MTTF of a single item!

Finally, the failure rate is given by

$$h_{S}(t) = -\frac{R_{S}'(t)}{R_{S}(t)} = \frac{6\lambda e^{-2\lambda t} - 6\lambda e^{-3\lambda t}}{3e^{-2\lambda t} - 2e^{-3\lambda t}} = \frac{6\lambda e^{-2\lambda t} (1 - e^{-\lambda t})}{e^{-2\lambda t} (3 - 2e^{-\lambda t})} = \frac{6\lambda \cdot \frac{1 - e^{-\lambda t}}{3 - 2e^{-\lambda t}}}{3 - 2e^{-\lambda t}}.$$

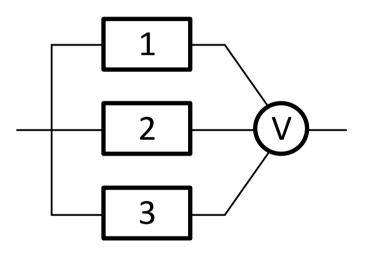


The *2-out-of-3* system is the most common case of *k-out-of-n* redundancy, since all other variants are more costly.

Often, 2-out-of-3 system is called TMR system, where TMR stands for triple modular redundancy.

Also, there is a subclass of *k*-out-of-*n* systems called <u>majority</u> <u>voting systems</u>. Here *n* is always odd number, and $k = \frac{n+1}{2}$.

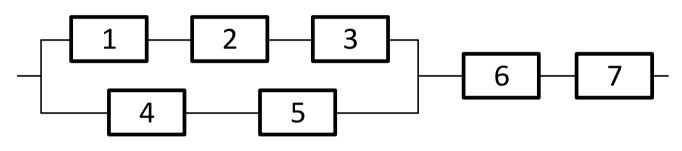
The RBD for 2-out-of-3 system is as follows:



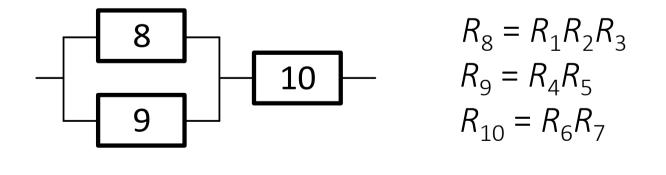
Often, the element V (voter) is assumed to be perfectly reliable, i.e. $R_V(t)=1$.

The formulae for the reliability computation of series and parallel systems can be used in combination to compute the reliability of a system having both series and parallel parts (*series-parallel systems*).

The computational procedure consists of a progressive reduction of the system complexity by substituting blocks of components in series/parallel with a single equivalent block. *Ex.*: Consider a system with the following RBD:



1st step: replace all series elements with equivalent ones:

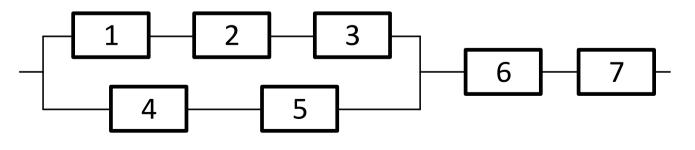


2nd step: replace parallel elements 8 and 9 with equivalent one:

$$-11 - 10 - R_{11} = R_8 + R_9 - R_8 R_9$$

Finally, reliability of the entire system is given by

$$R_S = R_{10}R_{11} = (R_1R_2R_3 + R_4R_5 - R_1R_2R_3R_4R_5) \cdot R_6R_7$$

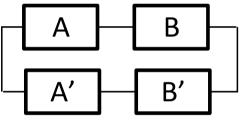


When discussing series-parallel systems, we should address another topic, namely, <u>system redundancy vs. component</u> <u>redundancy</u>.

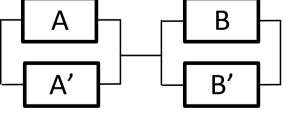
Consider a system composed of two series components A and B:

Let's denote its reliability as $R_1 = R_A R_B$.

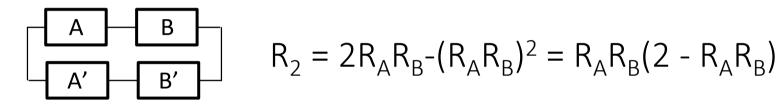
If we decide to improve the reliability of the system by applying redundancy using one single replica for each component, two solutions are possible. Either we replicate the complete line (system redundancy):



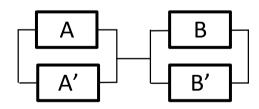
or we replicate each component individually (<u>component</u> <u>redundancy</u>):



Reliability of the system with system redundancy is



Reliability of the system with component redundancy:



$$R_{3} = (2R_{A} - R_{A}^{2})(2R_{B} - R_{B}^{2}) =$$
$$= R_{A}R_{B}[4 - 2(R_{A} + R_{B}) + R_{A}R_{B}]$$

 $R_{1} = R_{A}R_{B}$ $R_{2} = R_{A}R_{B}(2 - R_{A}R_{B})$ $R_{3} = R_{A}R_{B}[4 - 2(R_{A} + R_{B}) + R_{A}R_{B}]$

It is easy to see that $R_2 > R_1$ and $R_3 > R_1$; both redundant configurations are more reliable than the original system.

Next, we compare between the two redundant configurations

$$\frac{R_3}{R_2} = \frac{4 - 2(R_A + R_B) + R_A R_B}{2 - R_A R_B} = 1 + \frac{2(1 - R_A)(1 - R_B)}{2 - R_A R_B} > 1$$

It should, however, be noted that configuration 3 is more complex than configuration 2, since each type A component need to be possibly connected with any type B component.

This higher complexity requires an additional control logic (not considered in the formulae) that may reduce the benefits calculated from the equation.

The reason why configuration 3 is more reliable than configuration 2 can be also explained on a qualitative basis, noticing that there are failure combinations of basic blocks that cause failure of configuration 2, but not of configuration 3.

